

# Novel Sorting Algorithm

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**Abstract---** Sorting has become a stand in need of prosaic life activities. Sorting in computer science alluded to as ordering deals with arranging elements of a list or a set of records of a file in ascending or descending order. We have gone through a number of sorting algorithms like bubble sort, selection sort, insertion sort, shell sort, merge sort and bucket etc. Each algorithm has intrinsic odds and restrictions. We cannot say that a particular algorithm is the best algorithm, because algorithm may be easy to implement but it may take maximum time to execute, where as the other algorithm may be hard to implement but it may save execution time. The time complexity of sorting algorithms is, maximum in the range  $O(n)$  to  $O(n^2)$ . Most of the algorithms have  $O(n^2)$  as worst case time complexity and for few algorithms have  $O(n)$  as the best case time complexity. Some algorithms have  $O(n \log n)$  as the time complexity in best and or average and or worst cases.

We are proposing novel sorting algorithm which has time complexity  $O(n)$  in the best case and  $O(n^2)$  in the worst case.

**Keywords-** Time Complexity, Sorting, Best case, Worst case.

## I. INTRODUCTION

Sorting in English language refers to separating or arranging things according to different classes. Sorting in computer science, refers arranging data either in decreasing or increasing order. Sorting has become a stand in need of our daily life activities. For example a school going student has to align his/her class notes according to the time table. A job holder has to categorize his/her files according to the priorities of the file. A house maker has to arrange all her domestic commodity in a proper order for convenient use. Therefore sorting is playing a major role in our daily life.

Sorting organize the data in proper order. If the data is in proper order we can easily access the required data. If not, it is tough to access it because; we do not know where we can find the data. To find data we have to perform search operation on the entire database. If the data is in the order, searching successive elements is easy. If data is not sorted out then it is difficult to access successive elements because the elements are located at various locations.

In bubble sort algorithm [13], two elements are compared and rearranged if necessary. In the proposed method we consider three elements and move one element towards left or right and other element is moved in opposite direction. By using this we can minimize the number of swap operations and or iterations.

## II. PROPOSED NOVEL SORTING ALGORITHM

We studied behavior of elements when sorting method processed on the list of elements in most commonly used sorting algorithm along with time and space complexities of algorithms. Bubble sort and selection sorts have  $n(n-1)/2$  comparisons. The insertion sort has  $n$  comparisons in the best case where  $n$  is the number of elements in the list.

In our novel sorting algorithm, in the each iteration bigger element moved towards right like bubble sort and smaller element moved one or two positions towards left where as in the bubble sort only one element moved either direction only. This is the basic difference between bubble sort and the proposed algorithm.

To sort the given list of elements, we append largest element at the end of the list. This step is to generalize the coding part but it may or may not be useful as it depends on the number of elements. The appended element is useful only when the list contains even number of elements. In this algorithm, to minimize the swap operations at a time we compare element in the odd location (index 1,3,5...where index start from 0.) with the neighboring elements and arranged in the proper order.

For example consider the elements 8 5 3 9 1

➤ **First iteration (Bubble sort)**

Compare 8 and 5 then swap these elements, after swap operation the elements are 5 8 3 9 1  
 Compare 8 and 3 then swap these elements, after swap operation the elements are 5 3 8 9 1  
 Compare 8 and 9 but no swap because these elements are in order.  
 Compare 9 and 1 then swap these elements, after swap operation the elements are 5 3 8 1 9.

➤ **Second iteration (Novel sort)**

Consider list of element given in the above example 8 5 3 9 1  
 Element in first odd location is 5 and this element is compared with 8 and 3. Arrange the elements based on the order. Therefore the order is 3 5 8 9 1

Now element in the next odd location is 9 and this element compared with 8 and 1 (adjacent elements). The order is 3 5 1 8 9

After first iteration the result of Bubble sort and Novel is

Bubble sort 5 3 8 1 9

Novel sorting : 3 5 1 8 9

In proposed algorithm smaller elements also moved towards left (one or two positions) in each iteration.

➤ **Second iteration (Bubble sort)**

5 3 8 1 9

Compare 5 and 3, swap these two, after swap operation the elements are 3 5 8 1 9

Compare 5 and 8 but no swap, as they are in order

Compare 8 and 1, swap these two, after swap operation the elements are 3 5 1 8 9

➤ **Second iteration (Novel sort)**

3 5 1 8 9

Element 5 compared with 3 and 1 arrange elements 1 3 5 8 9

Bubble sort requires few more iterations to sort elements. The proposed algorithm does not require any more iteration. The advantage in the each iteration is, bigger elements are moved towards right and smaller elements are moved towards left to minimize the number of swap operations and or iterations.

**A. Novel sorting algorithm**

Input: list of elements  $a[0..n-1]$ , where  $n$  is number of elements

Step 1:  $m=n$ ,  $a[n]=\text{maximum}$

Step 2: repeat steps 3 and 5 for  $j=0$  to  $n/2$  where step size=1

Step 3: repeat step 4 for  $i=1$  to  $m$  where step size=2

Step 4: compare elements  $a[i-1]$ ,  $a[i]$  and  $a[i+1]$ . If they are not in order arrange them in order.

Step 5:  $m--$

The above algorithm sorts the elements. In some cases even though the elements are sorted in order the loop is unnecessarily executed. To overcome this drawback we can use flag variable for checking swap operation. If swap operations are not performed in inner loop then the elements are in the sorted order.

**B. Modified Novel sorting Algorithm**

Input: list of elements  $a[0..n-1]$ , where  $n$  is number of elements

Step 1:  $m=n$ ,  $\text{swap}=0$ ,  $a[n]=\text{maximum}$

Step 2: repeat step 3, 5, and 6 for  $j=0$  to  $n/2$  where step size=1

Step 3: repeat step 4 for  $i=1$  to  $m$  where step size=2

Step 4: compare elements  $a[i-1]$ ,  $a[i]$  and  $a[i+1]$ . If they are not in order arrange them in order.

Set  $\text{swap}=1$ ;

Step 5: if  $\text{swap}=0$  then given elements are in order break the outer loop else set  $\text{swap}=0$

Step 6:  $m--$

### III. ALGORITHM ANALYSIS

#### A. Best case:

If the elements are in order, then outer loop will be terminated after completion of the first iteration of inner loop. Therefore time complexity is  $O(n)$

#### B. Worst case:

We assume that list contain odd number of elements. The outer loop repeats for  $n/2$  times. In the first pass of outer loop, the inner loop repeats for  $(n/2)$  times and performs  $3n/2$  comparison operations. The number of assignments performed depends on the order of elements. The maximum number of assignments performed is  $2n$ .

In the second pass  $(n/2)$ , in third pass  $(n/2-1)$ .. .. .

Inner loop repeats  $n/2+n/2+n/2-1+n/2-1+n/2-2+n/2-2$ ..... $(n/2+1)$  terms.

i.e  $n/2+n/2+n/2-1+n/2-1+n/2-2+n/2-2$ ..... $(n/2+1)$  terms

$$\approx 7n^2/32 + 7n/8 \quad \text{when } n \text{ is very large then}$$

$$\approx n^2$$

Therefore the time complexity is  $O(n^2)$

### IV. RESULTS

The time complexity of this algorithm in good case is  $O(n)$  and in worst case  $O(n^2)$  same as bubble sort but their actual run time differ. To better understanding the actual performance we conducted some experiments. The run times are measured on a PC, AMD Athlon64X2 dual core 4200+ processor and 1G.B. RAM under Microsoft XP operating system. These algorithms are compiled using the sun java platform compiler and run under the java interpreter. The run time shown is CPU execution time measured using object of Date class. The class Date available in java util package. The elements are generated using nextInt method of Random class. The same set of elements is used for both algorithms.

Table 1. Execution time of Bubble and Novel Sorting Algorithms

N =	1000	5000	10000	20000	30000	40000	50000
<b>Bubble sort</b>	4.42	90.3	359.68	1457.76	3525.34	6536.56	10311.56
<b>Novel sorting method</b>	17.44	61.8	199.12	770.04	1791.92	3104.36	4859.96

### V. CONCLUSION

We have proposed a novel algorithm to sort given elements. The new algorithm compares three elements at a time and rearranges these elements. The proposed algorithm is easy to understand and easy to implement. The proposed novel algorithm has a similarity with bubble sort that is in every phase one element moved to its correct location.

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